

Direct Search Based on Probabilistic Descent

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University of Coimbra, moving to CERFACS-IRIT joint lab

Joint work with S. Gratton, C. W. Royer, and L. N. Vicente

SIOPT — May 22, 2014, San Diego

It is a privilege to conclude the whole conference!

Unconstrained derivative-free optimization (DFO)

$$\min_{x \in \mathbb{R}^n} f(x)$$

$$f : \mathbb{R}^n \rightarrow \mathbb{R}$$

f is **bounded from below** and **differentiable**
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- S. Gratton, P. Laloyaux, and A. Sartenaer, “Derivative-free Optimization for Large-scale Nonlinear Data Assimilation Problems”, 2013.

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And convergence **theory** of the algorithms.

- Two main classes of rigorous methods in DFO
 - Directional methods, like direct search (GPS, GSS, MADS ...)
 - Model-based methods, like trust region methods (DFO, NEWUOA, CONDER, BOOSTER, ORBIT ...)

Direct search (DS)

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For $k = 0, 1, 2, \dots$

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$$f(x_k + \alpha_k d_k) < f(x_k) - \rho(\alpha_k).$$

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$$x_{k+1} = \begin{cases} x_k + \alpha_k d_k & \text{if successful} \\ x_k & \text{if unsuccessful,} \end{cases}$$

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A **forcing function** ρ is a positive and monotonically **nondecreasing** function such that

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In this talk:

$$\rho(\alpha) = \frac{\alpha^2}{2}$$

$$\alpha_0 = 1 \quad (\text{initial stepsize})$$

$$\gamma = 2 \quad (\text{increasing factor})$$

$$\theta = \frac{1}{2} \quad (\text{decreasing factor})$$

Traditional polling set: PSS

- Positive spanning set (PSS):

$D = \{d_1, \dots, d_m\}$ is a PSS if it spans \mathbb{R}^n positively:

$$\mathbb{R}^n = \left\{ \sum_{i=1}^m \mu_i d_i : \mu_i \geq 0 \ (1 \leq i \leq m) \right\}.$$

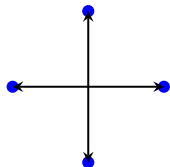
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Example:



$$D_{\oplus} = \{e_1, \dots, e_n, -e_1, \dots, -e_n\}$$

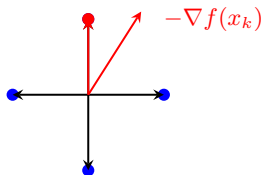
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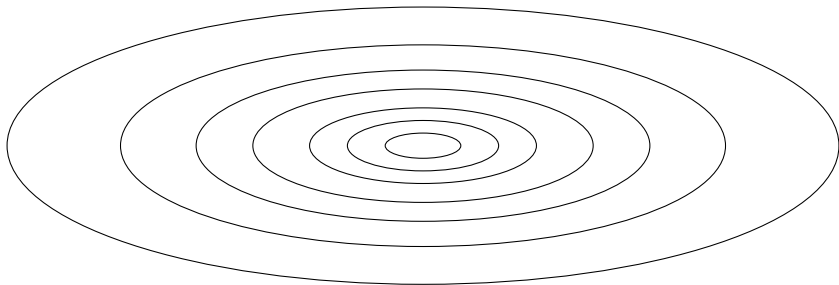
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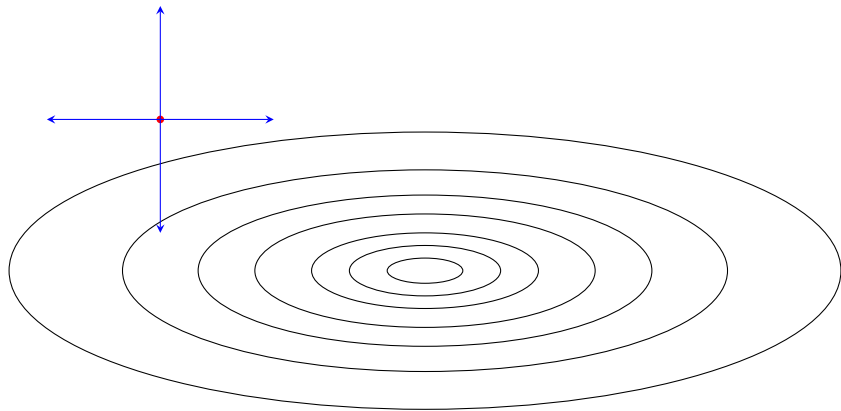
- $\exists d \in D$ that 'approximates' $-\nabla f(x_k)$, meaning $d^{\top}[-\nabla f(x_k)] > 0$.

Example: Coordinate search



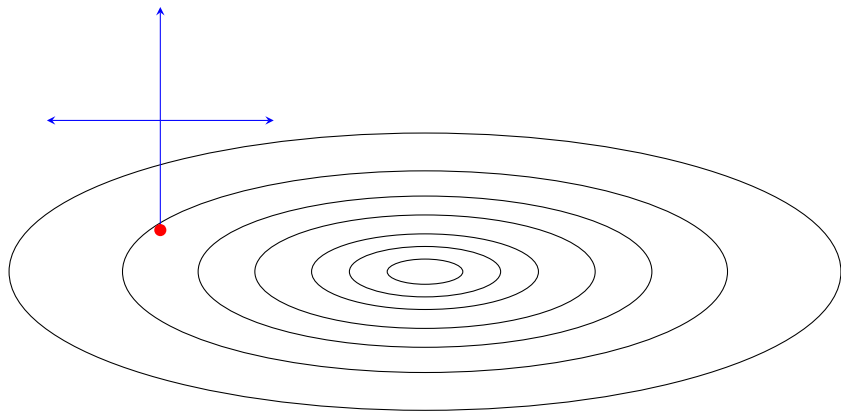
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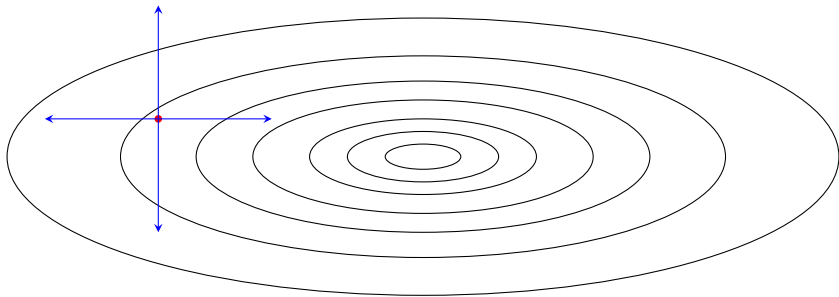
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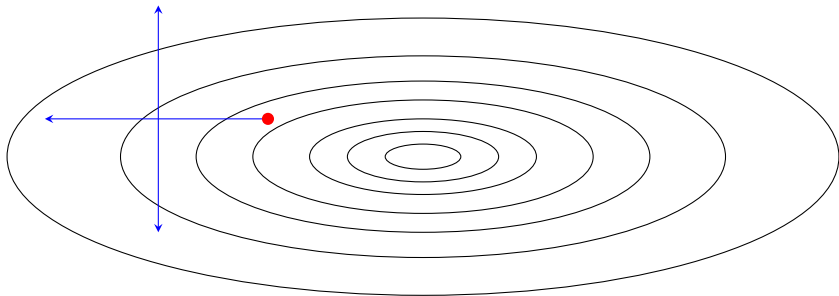
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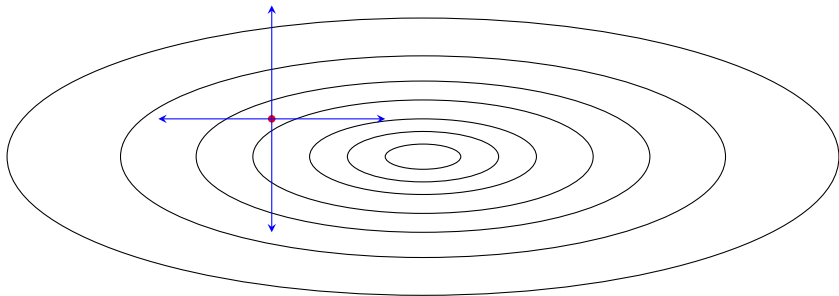
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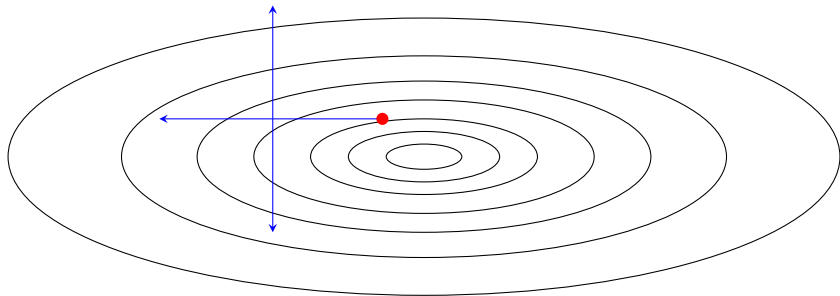
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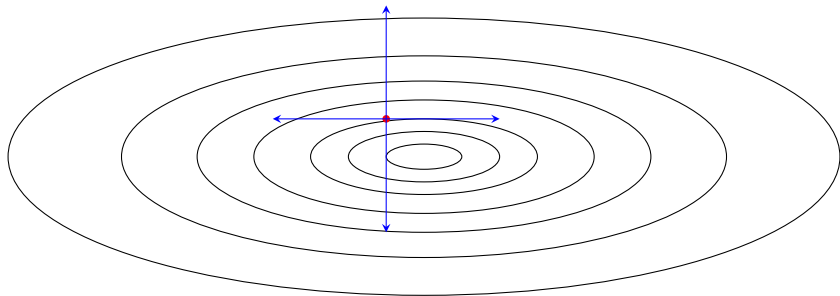
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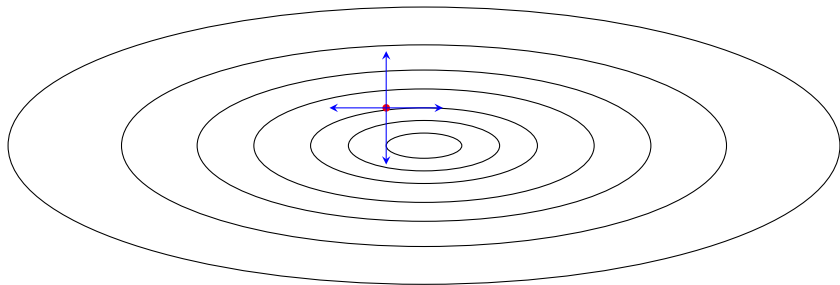
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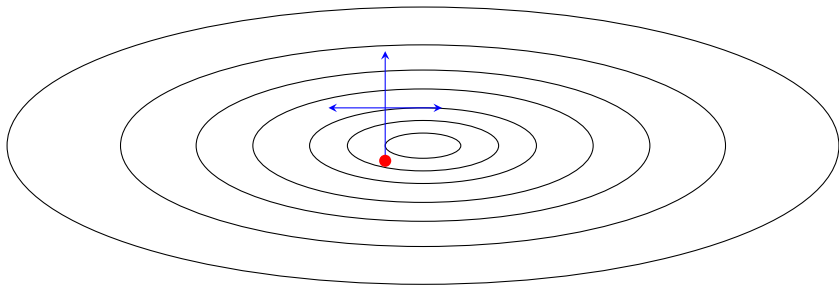
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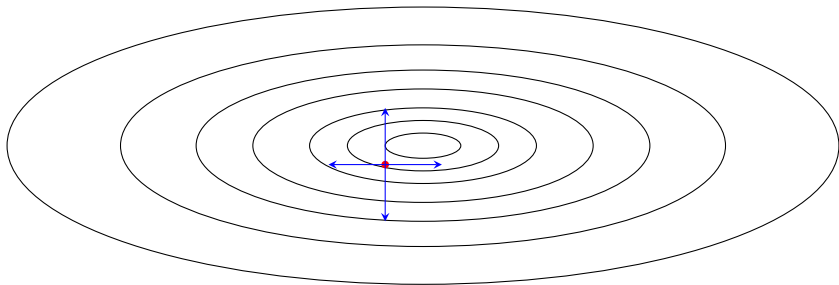
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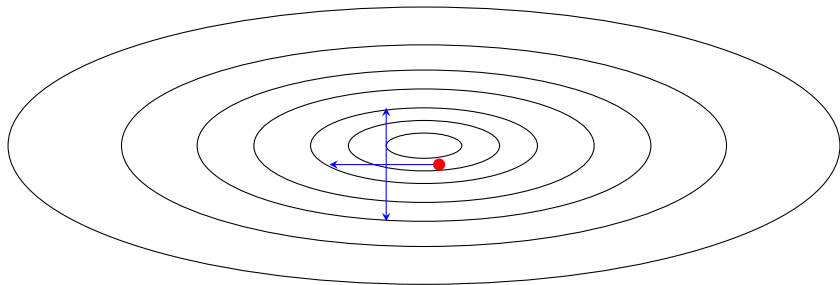
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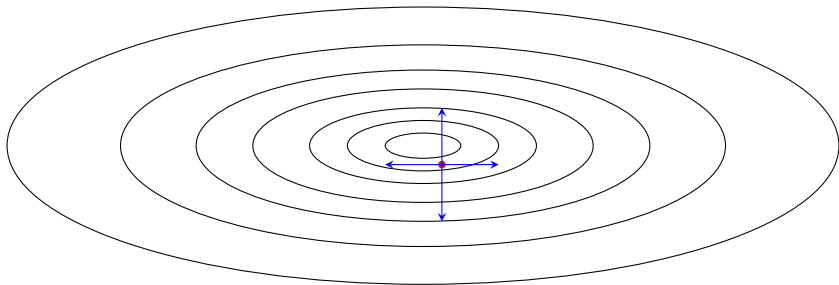
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- Cosine measure: the ability of D to 'approximate' directions in \mathbb{R}^n .

$$\text{cm}(D) = \min_{0 \neq v \in \mathbb{R}^n} \max_{d \in D} \frac{d^\top v}{\|d\| \|v\|}.$$

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- Example:

$$\text{cm}(D_\oplus) = \frac{1}{\sqrt{n}}.$$

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Theorem (Torczon 1997, Kolda, Lewis, and Torczon 2003)

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- $\min_{0 \leq \ell \leq k} \|\nabla f(x_\ell)\| = \mathcal{O}(1/\sqrt{k})$.
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DS with PSS: Theory

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Question: Does the theory cover the most efficient implementation of DS?

A competitor against PSS: Random polling set

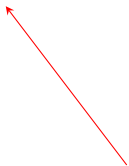
- Success of [random coordinate descent](#), [stochastic gradient](#) ...
 - Y. Nesterov, “Efficiency of coordinate descent methods on huge-scale optimization problems”, SIAM Journal on Optimization, 22(2), 341-362, 2012
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- Success of randomization in derivative-free methods, with insightful theories:
 - A. S. Bandeira, K. Scheinberg, and L. N. Vicente, “Convergence of trust-region methods based on probabilistic models”, submitted
 - K. Scheinberg, “Convergence rates of line-search and trust region methods based on probabilistic models”, MS16, SIOPT14

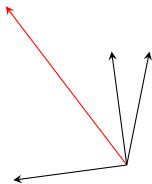
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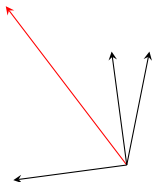


$n + 1$ random polling directions

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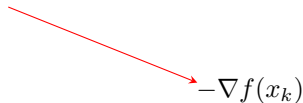
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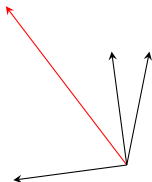
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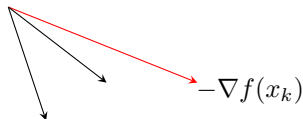
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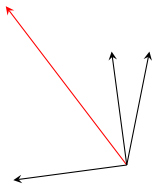


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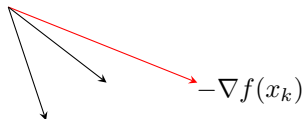
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D_k is 'good' in some probabilistic sense ...

What do we mean by 'good'?

If derivatives were available, it would have been sufficient to require

$$\max_{d \in D} \frac{-d^\top \nabla f(x_k)}{\|d\| \|\nabla f(x_k)\|} \geq \kappa.$$

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Then $\text{cm}(D, -\nabla f(x_k)) \geq \kappa$ would have been enough.

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But derivatives are not available!

Random variables v.s. realizations

From now on, we suppose that the polling directions are **not defined deterministically** but taken at **random** from the **unit sphere** \mathcal{S}^{n-1} .

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Distinguish **random variables** from **realizations**

	Iterate	Polling set
Random variables	X_k	\mathcal{D}_k
Realizations	x_k	D_k

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- Global rate:

$$\mathbb{P} \left(\min_{0 \leq \ell \leq k} \|\nabla f(X_k)\| \leq \frac{C}{\sqrt{k}} \right) \text{ is 'high'}$$

for some properly selected constant C ?

How to achieve the goals?

- Global convergence:

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$$\left\{ \min_{0 \leq \ell \leq k} \|\nabla f(X_k)\| > \epsilon \right\} \subset E_{k,\epsilon},$$

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Let us find E and $E_{k,\epsilon}$...

Global convergence: An intuitive lemma

Let Z_k be the indicator function of $\{\text{cm}(\mathfrak{D}_k, -\nabla f(X_k)) \geq \kappa\}$, and

$$p_0 = \frac{\ln \theta}{\ln(\gamma^{-1}\theta)} = \frac{1}{2}.$$

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Without imposing any assumption on the probabilistic behavior of $\{\mathfrak{D}_k\}$:

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$$\left\{ \liminf_{k \rightarrow \infty} \|\nabla f(X_k)\| > 0 \right\} \subset \left\{ \sum_{k=0}^{\infty} (Z_k - p_0) = -\infty \right\}.$$

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Meaning:

If convergence does not hold, the ‘frequency’ of $\{Z_k\}_{k \geq 0}$ is ‘less than p_0 ’.

Global rate: Another intuitive lemma

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Lemma

$$\left\{ \max_{0 \leq \ell \leq k} \|\nabla f(X_k)\| > \epsilon \right\} \subset \left\{ \sum_{\ell=0}^{k-1} Z_\ell \leq \left[\frac{(L+1)^2 \beta}{2\kappa^2 \epsilon^2 k} + p_0 \right] k \right\}.$$

$\beta < \infty$ is an upper bound for $\sum_{k=0}^{\infty} \rho(\alpha_k)$ (existence guaranteed).

$L < \infty$ is a Lipschitz constant of ∇f in \mathbb{R}^n .

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Lemma

$$\left\{ \max_{0 \leq \ell \leq k} \|\nabla f(X_k)\| > \epsilon \right\} \subset \left\{ \sum_{\ell=0}^{k-1} Z_\ell \leq \left[\frac{(L+1)^2 \beta}{2\kappa^2 \epsilon^2 k} + p_0 \right] k \right\}.$$

$\beta < \infty$ is an upper bound for $\sum_{k=0}^{\infty} \rho(\alpha_k)$ (existence guaranteed).

$L < \infty$ is a Lipschitz constant of ∇f in \mathbb{R}^n .

Meaning:

If $\{\|\nabla f(X_0)\|\}_{0 \leq \ell \leq k}$ are all above ϵ , the 'frequency' of $\{Z_\ell\}_{0 \leq \ell \leq k-1}$ is 'not more than' $p_0 + \mathcal{O}(\epsilon^{-2} k^{-1})$.

What assumptions to impose?

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Definition

The sequence $\{\mathfrak{D}_k\}$ is *p -probabilistically κ -descent* if, for each $k \geq 0$,

$$\mathbb{P}(\text{cm}(\mathfrak{D}_k, -\nabla f(X_k)) \geq \kappa \mid \mathfrak{D}_0, \dots, \mathfrak{D}_{k-1}) \geq p.$$

Lemma

If $\{\mathfrak{D}_k\}$ is p_0 -probabilistically κ -descent, then $\left\{ \sum_{\ell=0}^{k-1} (Z_\ell - p_0) \right\}$ is a submartingale, and

$$\mathbb{P} \left(\sum_{k=0}^{\infty} (Z_k - p_0) = -\infty \right) = 0.$$

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Theorem

If $\{\mathfrak{D}_k\}$ is p_0 -probabilistically κ -descent, then

$$\mathbb{P} \left(\liminf_{k \rightarrow \infty} \|\nabla f(X_k)\| = 0 \right) = 1.$$

The analysis is inspired by that for probabilistic trust region.

Lemma (Chernoff bound)

Suppose that $\{\mathfrak{D}_k\}$ is *p -probabilistically κ -descent* and $\lambda \in (0, p)$. Then

$$\mathbb{P} \left(\sum_{\ell=0}^{k-1} Z_{\ell} \leq \lambda k \right) \leq \exp \left[-\frac{(p - \lambda)^2}{2p} k \right].$$

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Theorem

Suppose that $\{\mathfrak{D}_k\}$ is p -probabilistically κ -descent with $p > p_0$. Then

$$\mathbb{P} \left(\min_{0 \leq \ell \leq k} \|\nabla f(X_{\ell})\| \leq \left[\frac{(L+1)\beta^{\frac{1}{2}}}{(p-p_0)^{\frac{1}{2}}\kappa} \right] \frac{1}{\sqrt{k}} \right) \geq 1 - \exp \left[-\frac{(p-p_0)^2}{8p} k \right].$$

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$\implies \mathcal{O}(1/\sqrt{k})$ decaying rate for gradient holds with overwhelmingly high probability, matching the deterministic case (Vicente 2013).

Practical probabilistic descent sets

For each $k \geq 0$,

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For each $k \geq 0$,

- \mathfrak{D}_k is independent of the previous iterations,
- \mathfrak{D}_k is a set $\{\mathfrak{d}_1, \dots, \mathfrak{d}_m\}$ of independent random vectors uniformly distributed on the unit sphere.

Practical probabilistic descent sets

$\{\mathfrak{D}_k\}$ generated in this way is probabilistically descent.

Proposition

Given $\tau \in [0, \sqrt{n}]$, $\{\mathfrak{D}_k\}$ is p -probabilistically (τ/\sqrt{n}) -descent with

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For instance,

$$\left. \begin{array}{l} m = 2 \\ \tau = \frac{1}{2} \end{array} \right\} \implies p > \frac{1}{2} = p_0.$$

Practical probabilistic descent sets: WCC bounds

Plugging $\kappa = 1/(2\sqrt{n})$ into the global rate, one obtains

WCC (number of iterations)

$$\mathbb{P} \left(K_{\epsilon} \leq \left\lceil \frac{4(L+1)^2\beta}{p-p_0} (n\epsilon^{-2}) \right\rceil \right) \geq 1 - \exp \left[- \frac{\beta(p-p_0)(L+1)^2}{2p} (n\epsilon^{-2}) \right].$$

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$$\mathbb{P} \left(K_{\epsilon}^f \leq 2 \left\lceil \frac{4(L+1)^2 \beta}{p - p_0} (n\epsilon^{-2}) \right\rceil \right) \geq 1 - \text{the tiny tail.}$$

$\implies \mathcal{O}(n\epsilon^{-2})$ with overwhelmingly high probability, better than the deterministic case $\mathcal{O}(n^2\epsilon^{-2})$ (Vicente 2013).

The competition

Relative performance: PSS v.s. Random polling sets ($n = 40$)

	D_{\oplus}	$2n$	$n + 1$	$n/4$	2	1
arglina	3.42	10.30	6.01	1.88	1.00	—
arglinb	20.50	7.38	2.81	1.85	1.00	2.04
broydn3d	4.33	6.54	3.59	1.28	1.00	—
dqrtic	7.16	9.10	4.56	1.70	1.00	—
engval1	10.53	11.90	6.48	2.08	1.00	2.08
freuroth	56.00	1.00	1.67	1.67	1.00	4.00
integreq	16.04	12.44	6.76	2.04	1.00	—
nondquar	6.90	7.56	4.23	1.87	1.00	—
sinqquad	—	1.65	2.01	1.00	1.55	—
vardim	1.00	1.80	2.40	1.80	1.80	4.30

Solution accuracy was 10^{-3} . Averages were taken over 10 independent runs.

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- An interesting future work: randomized subspace method.

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- Probabilistic DS enjoys, with overwhelmingly high probability:
 - the same WCC for number of iterations,
 - **possibly better** WCC for number of function evaluations.
- The analysis technique can be applied to **probabilistic trust region** method $\implies \mathcal{O}(1/\sqrt{k})$ rate for gradient.
- An interesting future work: randomized subspace method.
Let $\text{Gr}(l, \mathbb{R}^n)$ be the set of all the l -dim linear subspaces of \mathbb{R}^n .

Lemma (Randomized subspace)

Suppose that S is **uniformly** distributed on $\text{Gr}(l, \mathbb{R}^n)$. Then for any nonzero vector $v \in \mathbb{R}^n$ and constant $\delta \in (0, 1)$,

$$\mathbb{P} \left(\|Pv\| \geq \sqrt{\frac{l\delta}{n}} \|v\| \right) \geq 1 - \exp \left[-\frac{l}{2} (\delta - 1 - \ln \delta) \right].$$